IN THE SPECIFICATION

Please amend the paragraph at page 42, lines 4-15, as follows:

Dth is, for example, ~~(k + 1)~~(k + ℓ). Here k is such a number that the product of N(1), …, N(k) is not less than the total number of user IDs (identification information items) (N(1) ≦ N(2) ≦ … ≦ N(M)), and [[l]] ℓ is given by the following formula (1):

$$[1-\Pi\ 1/N(i)]^S \geqq 1 - \varepsilon_2 \qquad (1)$$

where the range of i that assumes ~~Π is i = 1 ~ l~~ Π is i = 1 ~ ℓ or ~~i = k+1 ~ (k + 1)~~ i = k+1 ~ (k + ℓ),

~~S = $_M C_{k+1}$~~ S = $_M C_{k+\ell}$, and

ε_2 represents the rate of error tracing in each user ID of the people responsible for collusive attacks, and satisfies $0 < \varepsilon_2 < 1$.

Please amend the paragraph beginning at page 52, line 27, to page 53, line 1, as follows:

Assume that ~~M = c (k + 1)~~ M = c (k + ℓ), C is a narrow sense $[M, k, M-k+1]_q$ Reed-Solomon code.

Please amend the paragraph at page 53, lines 2-11, as follows:

If the following formula (2) is satisfied, the Reed-Solomon code C can be made to be a stochastic outer code:

~~$[1 - 1/q^1]^S \geqq 1 - \epsilon$~~ $[1 - 1/q^{\ell}]^S \geqq 1 - \varepsilon$ (2)

where ~~S = $_M C_{k+1}$~~ S = $_M C_{k+\ell}$,

$q = N(1) = N(2) = \ldots = N(M)$,

and ε represents the rate of error tracing in each user ID (identification information) of the people responsible for collusive attacks, and is a real number that satisfies $0 < \varepsilon < 1$.

Please amend the paragraph at page 53, lines 12-17, as follows:

In this case, the above-described tracing algorithm example as a stochastic method is applicable. In the tracing algorithm example as a stochastic method, [[l]] ℓ included in the formula, ~~Dth = k + l~~ Dth = k + ℓ, may be given by, for example, formula (2) instead of formula (1).